

# Second XCSP3 Competition (2018 CSP and COP competition)

Last Call for Solvers and Benchmarks  
(**Strict Participation Deadline:** April 29, 2018)

<http://www.xcsp.org>

The second international XCSP3 constraint solver competition is organized to improve our knowledge about components (e.g., filtering algorithms, heuristics, search strategies, and learning procedures) that are behind the efficiency of solving systems (referred to as constraint solvers in this document) for combinatorial constrained problems. Two classical problems are considered for this competition:

- CSP (Constraint Satisfaction Problem)
- COP (Constrained Optimization Problem)

The intermediate<sup>1</sup> format XCSP3 is used as input format for the solvers. The effort required for entering the competition is limited because some tools (parsers) are available, and only a central set of popular (and important) constraints is considered.

This call for solvers and benchmarks presents the tracks that will be considered during the competition. In particular, we give important details about the format restrictions, the execution environment, and the rules that must be followed by the solvers.

Importantly, do note that:

- you can find the final and detailed results concerning the first XCSP3 competition on [www.cril.fr/XCSP17](http://www.cril.fr/XCSP17). Nine teams were competing in 2017 for the first edition.
- anybody (an in particular, contestants) is strongly solicited to submit new benchmarks (you can, for example, use the Java-based API called [MCSP3](#) for that).

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<sup>1</sup>XCSP3 is neither a modeling language, nor a flat format. It is intermediate because it preserves the structure of problems through the concepts of variable arrays, constraint groups/blocks, and meta-constraints.

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## 1 Timetable

The deadlines of the competition are defined below:

Opening of the registration site at <a href="http://www.cril.fr/XCSP18/">http://www.cril.fr/XCSP18/</a>	March 2018
Pre-registration of contestants	March 20, 2018
Final registration (submission of solvers and benchmarks)	April 29, 2018
Test of solvers conformance	early May 2018
Position paper (2 pages)	mid-June 2018
Competition running	May-June 2018
Final results available	during CP 2018

Once submitted, solvers will be run on a limited number of benchmarks to make sure that they interact correctly with the evaluation environment. Potential problems will be reported to the authors by the 13th of May 2018. Bug fixes will be due by the 20th of May 2018.

## 2 Tracks

First, do note that we consider two main problems: CSP (a decision problem) and COP (an optimization problem). For CSP, the goal is to exhibit one solution or to prove that none exists. For COP (mono-objective optimization), the goal is to exhibit a solution with the best possible objective value, ideally proving that it represents an optimum solution.

Anyone can submit a solver to any particular track. There are exactly 8 tracks. The first 6 tracks impose absolutely no conditions on solvers. For example, they can be written in any language (provided that we can reasonably execute them in our environment), and can be complete or incomplete solvers (e.g., based on local search). The 6 Standard tracks are described by Table 1.

<b>Problem</b>	<b>Goal</b>	<b>Exploration</b>	<b>Timeout</b>
CSP	one solution	sequential	40 minutes
CSP	one solution	parallel	40 minutes
COP	best solution	sequential	4 minutes
COP	best solution	parallel	4 minutes
COP	best solution	sequential	40 minutes
COP	best solution	parallel	40 minutes

Table 1: Standard Tracks.

Ranking for COP will be stated in two different manners: by considering and not considering possible proofs of optimality, permitting in the latter case to emphasize the quality of incomplete solvers.

There are also 2 Mini-Solver tracks, where a mini-solver is a solver whose code must be open-source (and expected to be rather compact, although we do not impose anymore specific restrictions). For Mini-solvers tracks, only sequential exploration is considered, and the set of constraints is rather restricted, as described in the next section.

Our interest in Mini-solver tracks is two-fold. First, this should facilitate the participation of people (e.g., students) whose solvers cannot be complete enough to compete with well-established solvers from the community. Second, it will hopefully provide the community with one (or several) compact well-documented and easily extendable open-source solver.

<b>Problem</b>	<b>Goal</b>	<b>Exploration</b>	<b>Timeout</b>
CSP	one solution	sequential	40 minutes
COP	best solution	sequential	40 minutes

Table 2: Mini-Solver Tracks.

### 3 Format

The complete description of the format (XCSP3) used to represent combinatorial constrained problems can be found in [BLAP17]. Do note that we refer to the **version 3.0.5** of the specifications. However, for the 2018 competition, we limit XCSP3 to its kernel, called XCSP3-core. This means that the scope of XCSP3 is restricted to:

- integer variables,
- CSP and COP problems,
- a set of 21 popular (global) constraints for Standard tracks, and a small set of constraints for Mini-solver tracks.

For simplicity, we also impose the following restrictions:

- Integer variables have finite domains (and so, the special value *infinity* is forbidden).
- Variable arrays always start indexing at 0 (and so, the attribute `startIndex`, whose default value is 0, cannot be associated with `<array>`).
- The attribute `as` can only be associated with elements `<var>` and `<array>`; see Section 10.5 in [BLAP17].
- Undefined variables are not accepted but useless variables are (note that parsers/solvers can easily identify useless variables); see Section 2.10 in [BLAP17].
- Advanced forms of constraints (see Part III in [BLAP17]) are not accepted, except for the very specific cases explicitly described in the rest of this section.
- Reification is not accepted. View extensions are only accepted for constraints `allDifferent` and `sum`, as explained later.
- The type of the objective (in case of a COP instance) cannot be "product" or "lex", and when it is "expression", the content of the element `<minimize>` or `<maximize>` can only be a variable (identifier); see Chapter 3 in [BLAP17].
- Any integer value occurring in an XCSP3 file must belong to the interval  $-2^{31}..2^{31} - 1$ .

### 3.1 Constraints for Standard Tracks

In the Standard tracks, we find twenty-one constraints. In practice, it turns out that specific code of propagators is needed for approximately 14 constraints only, because:

- similar propagators may be used for `regular` and `mdd`,
- similar propagators may be used for `maximum` and `minimum`,
- `channel` can be decomposed into `element` constraints,
- `ordered`, `allEqual` and `instantiation` can be *trivially* reformulated as `intension`,
- `slide` is decomposed into a set of constraints `intension` or `extension`; the parser can do it automatically for you.

Do note that a large majority of the 23,000 instances (from around 90 classical models, some of them being described for example at [CSPLib](#)) that are currently available on our website only involve these 21 constraints.

#### 3.1.1 Constraint `intension`

This constraint is described in Section 4.1.1.1 in [\[BLAP17\]](#). There is no competition restriction for this constraint.

#### 3.1.2 Constraint `extension`

This constraint is described in Section 4.1.1.2 in [\[BLAP17\]](#). Competition restrictions:

1. Compressed tables (i.e., tables with compressed tuples) and smart tables are not accepted. However, do note that short tables (i.e., tables with tuples containing `'*`) are accepted in 2018.
2. Empty Tables (i.e., tables with with 0 support or 0 conflict) are not accepted.

Note that unary, binary and n-ary extensional constraints are accepted.

#### 3.1.3 Constraint `regular`

This constraint is described in Section 4.1.2.1 in [\[BLAP17\]](#). Competition restrictions:

1. The automaton on which is based the constraint must be deterministic.

#### 3.1.4 Constraint `mdd`

This constraint is described in Section 4.1.2.3 in [\[BLAP17\]](#). Competition restrictions:

1. There must be at least one path from the root node to the terminal node.
2. In `<transitions>`, the root note is given by the first item of the first transition.

### 3.1.5 Constraint `allDifferent`

This constraint is described in Section 4.1.3.1 in [BLAP17]. Competition restrictions:

1. If present, the element `<except>` only contains one (integer) value.
2. Restricted forms (obtained by using the attribute `restriction`) are not accepted.

In addition to the basic form of `allDifferent`, the advanced form `allDifferent-matrix` described in Section 7.2.1 in [BLAP17] is accepted.

Also, handling view extensions is authorized for the basic form of `allDifferent`. It means that instead of a list of variables inside the element `<list>`, it is possible to have a list of integer expressions (trees). This is shown at the end of Section 4.1.3.1 in [BLAP17]. Competition restriction:

3. For the basic form of `allDifferent`, the element `<list>` contains either only variables or only integer expressions (trees). It means that it is not possible to mix both forms.

### 3.1.6 Constraint `allEqual`

This constraint is described in Section 4.1.3.2 in [BLAP17]. There is no competition restriction for this constraint.

### 3.1.7 Constraint `ordered`

This constraint is described in Section 4.1.3.4 in [BLAP17]. Competition restrictions:

1. The compact form, obtained by using the attribute `case`, is not accepted.

Contrary to 2017, note that it is now possible to deal with an element `<lengths>`.

### 3.1.8 Constraint `lex`

This constraint is described in Section 7.1.4.1 in [BLAP17]. There is no competition restriction for this constraint.

In addition to this form of `lex`, the advanced form `lex-matrix` described in Section 7.2.2 in [BLAP17] is accepted.

### 3.1.9 Constraint `sum`

This constraint is described in Section 4.1.4.1 in [BLAP17]. Competition restrictions:

1. The condition is such that either the operator must be relational (i.e., must be in `{lt,le,gt,ge,eq,ne}`) and the (right) operand must be a value or a variable, or the operator must necessarily be `in` and the (right) operand must be an integer interval; See Section 1.5 in [BLAP17].

Also, handling view extensions is authorized for `sum`. It means here that instead of a list of variables inside the element `<list>`, it is possible to have a list of integer expressions (trees). This is shown at the end of Section 4.1.4.1 in [BLAP17]. Competition restriction:

2. The element `<list>` contains either only variables or only integer expressions (trees). It means that it is not possible to mix both forms.

### 3.1.10 Constraint count

This constraint is described in Section 4.1.4.2 in [BLAP17]. Competition restrictions:

1. The element `<values>` can only contain (integer) values.
2. The condition is such that either the operator must be relational (i.e., must be in `{lt,le,gt,ge,eq,ne}`) and the (right) operand must be a value or a variable, or the operator must necessarily be `in` and the (right) operand must be an integer interval; See Section 1.5 in [BLAP17].

### 3.1.11 Constraint nValues

This constraint is described in Section 4.1.4.3 in [BLAP17]. Competition restrictions:

1. If present, the element `<except>` only contains one (integer) value.
2. The condition is such that either the operator must be `eq` and the (right) operand must be a value or a variable, or the operator is `gt` and the (right) operand is 1.
3. Restricted forms (obtained by using the attribute `restriction`) are not accepted.

Contrary to 2017, note that it is now possible to deal with the special case where the condition is composed of the operator `gt` and the (right) operand is the value 1, which corresponds to the global constraint `notAllEqual`.

### 3.1.12 Constraint cardinality

This constraint is described in Section 4.1.4.4 in [BLAP17]. Competition restrictions:

1. The element `<values>` can only contain (integer) values.
2. Restricted forms (obtained by using the attribute `restriction`) are not accepted.

### 3.1.13 Constraint maximum

This constraint is described in Section 4.1.5.1 in [BLAP17]. Competition restrictions:

1. The condition is such that the operator must necessarily be `eq` and the (right) operand must be a value or a variable.
2. The element `<index>`, used for the variant `<arg_max>`, is not accepted.

### 3.1.14 Constraint minimum

This constraint is described in Section 4.1.5.2 in [BLAP17]. Competition restrictions:

1. The condition is such that the operator must necessarily be `eq` and the (right) operand must be a value or a variable.
2. The element `<index>`, used for the variant `<arg_min>`, is not accepted.



### 3.1.15 Constraint element

This constraint is described in Section 4.1.5.3 in [BLAP17]. Competition restrictions:

1. The optionnal attribute `startIndex`, if present, is necessarily equal to 0.
2. The attribute `rank`, whose default value is "any", cannot be present.

Contrary to 2017, it is now possible to have a list of values (instead of variables) inside `<list>`. This is presented in Section 4.1.5.3 in [BLAP17].

### 3.1.16 Constraint channel

This constraint is described in Section 4.1.5.4 in [BLAP17]. Competition restrictions:

1. Restricted forms (obtained by using the attribute `restriction`) are not accepted.

Contrary to 2017, note that for the form of `channel` involving two lists, it is now possible that these two lists have different sizes. This is discussed at the top of Page 79 in Section 4.1.5.4 in [BLAP17].

### 3.1.17 Constraint noOverlap

This constraint is described in Section 4.1.6.2 in [BLAP17]. Competition restrictions:

1. In case the element `<lengths>` contains (integer) values, whatever is the dimension, the value 0 is not accepted.

### 3.1.18 Constraint cumulative

This constraint is described in Section 4.1.6.3 in [BLAP17]. Competition restrictions:

1. The element `<ends>` is not accepted.
2. The variant, using `<machines>`, is not accepted.

### 3.1.19 Constraint instantiation

This constraint is described in Section 4.1.8.2 in [BLAP17]. There is no competition restriction for this constraint.

### 3.1.20 Constraint circuit

This constraint is described in Section 4.1.7.1 in [BLAP17]. Competition restrictions:

1. The optionnal attribute `startIndex`, if present, is necessarily equal to 0.
2. The element `<size>` cannot be present.

Contrary to 2017, we can deal with this constraint in 2018.

### 3.1.21 Meta-Constraint slide

This meta-constraint is described in Section 8.1 in [BLAP17]. Competition restrictions:

1. Only one element `<list>` is accepted.
2. The constraint template must be of form `<intension>` or `<extension>`.

## 3.2 Constraints for Mini-solver Tracks

All general restrictions introduced for Standard tracks hold. Additionnally, the constraints that are accepted for Mini-solver tracks are restricted to five types, as deccribed below.

### 3.2.1 Constraint intension

This constraint is described in Section 4.1.1.1 in [BLAP17]. Competition restrictions for Mini-solver Tracks : only primitive constraints with one of the following form will be considered. In what follows,  $x$ ,  $y$  and  $z$  denote integer variables,  $k$  denotes an integer value,  $\odot$  denotes a relational operator in  $\{<, \leq, \geq, >, =, \neq\}$  and  $\oplus$  denotes a binary arithmetic operator in  $\{+, -, *, /, \%, ||\}$ , with  $||$  being the distance.

- $x \odot k$                        $k \odot x$
- $x \odot y$
- $(x \oplus k) \odot y$        $(k \oplus x) \odot y$        $x \odot (y \oplus k)$        $x \odot (k \oplus y)$
- $(x \oplus y) \odot z$        $x \odot (y \oplus z)$

### 3.2.2 Constraint extension

This constraint is described in Section 4.1.1.2 in [BLAP17]. Competition restrictions for Mini-solver Tracks (the same as for Standard Tracks):

1. Compressed tables (i.e., tables with compressed tuples) and smart tables are not accepted. However, do note that short tables (i.e., tables with tuples containing ’\*’) are accepted in 2018.
2. Empty Tables (i.e., tables with with 0 support or 0 conflict) are not accepted.

Note that unary, binary and n-ary extensional constraints are accepted.

### 3.2.3 Constraint allDifferent

This constraint is described in Section 4.1.3.1 in [BLAP17]. Competition restrictions for Mini-solver Tracks:

1. The element `<except>` cannot be present.
2. Restricted forms (obtained by using the attribute `restriction`) are not accepted.

No advanced form of `allDifferent`, as e.g., `allDifferent-matrix`, is accepted.

Contrary to Standard tracks, handling view extensions for `allDifferent` is not permitted for the Mini-solver tracks.

### 3.2.4 Constraint sum

This constraint is described in Section 4.1.4.1 in [BLAP17]. Competition restrictions for Mini-solver Tracks:

1. The condition is such that the operator must be relational (i.e., in {**lt**, **le**, **gt**, **ge**, **eq**, **ne**}) and the (right) operand must be a value or a variable; see Section 1.5 in [BLAP17].

Contrary to Standard tracks, handling view extensions for **sum** is not permitted for the Mini-solver tracks.

### 3.2.5 Constraint element

This constraint is described in Section 4.1.5.3 in [BLAP17]. Competition restrictions for Mini-solver Tracks:

1. The optionnal attribute **startIndex**, if present, is necessarily equal to 0.
2. The element **<index>** is necessarily present but the attribute **rank**, whose default value is "any", cannot be present.

Contrary to 2017, it is now possible to have a list of values (instead of variables) inside **<list>**. This is presented in Section 4.1.5.3 in [BLAP17].

## 4 Differences between 2017 and 2018 Competitions

The differences in terms of restrictions, between the 2017 competition and the 2018 competition, are as follows:

1. Constraint **extension**. Contrary to 2017, short tables (i.e., tables with some tuples containing **'\***) are accepted in 2018. Note that an algorithm such as STR (Simple Tabular Reduction) or CT (Compact-Table) can be easily extended to deal with short tables.
2. Constraint **allDifferent**. Contrary to 2017, handling view extensions is authorized for the basic form of **allDifferent** in Standard tracks. It means that instead of a list of variables in the element **<list>**, it is possible to have a list of integer expressions (trees). This is shown at the end of Section 4.1.3.1 in [BLAP17].
3. Constraint **ordered**. Contrary to 2017, it is now possible to deal with an element **<lengths>**, as explained in Section 4.1.3.4 in [BLAP17].
4. Constraint **sum**. Contrary to 2017, handling view extensions is authorized for **sum** in Standard tracks. It means that instead of a list of variables in the element **<list>**, it is possible to have a list of integer expressions (trees). This is shown at the end of Section 4.1.4.1 in [BLAP17].
5. Constraint **nValues**. Contrary to 2017, note that it is now possible to deal with the special case where the condition is composed of the operator **gt** and the (right) operand is the value 1, which corresponds to the global constraint **notAllEqual**.

6. Constraint **element**. Contrary to 2017, it is now possible to have a list of values (instead of variables) inside `<list>`. See Section 4.1.5.3 in [BLAP17].
7. Constraint **channel**. Contrary to 2017, for the form of **channel** involving two lists, it is now possible that these two lists have different sizes. This is discussed at the top of Page 79 in Section 4.1.5.4 in [BLAP17].
8. Constraint **circuit**. This constraint is introduced in 2018. See restrictions in Section 3.
9. For some instances (series), the set of decision variables will be specified, by means of annotations.

## 5 Resources: Benchmarks and Tools

Many benchmarks can be found at:

[www.xcsp.org/series](http://www.xcsp.org/series)

The organizers invite *anybody* to submit new benchmarks. The organizers are particularly interested in new problem instances originating from real-world applications. For generating new XCSP3 instances, one can use the Java-based modeling API, called MCSP3. See its description in this [directory](#); once a model is developed, it is easy to generate XCSP3 instances by compiling it while providing data.

Some tools are also provided. They can be found at:

[www.xcsp.org/tools](http://www.xcsp.org/tools)

Currently, you can find:

- a C++ parser
- a Java parser
- a tool for checking solutions and costs
- a Java-based modeling API

## 6 Execution Environment

Solvers will run on a cluster of computers using the Linux operating system. They will run under the control of another program (called `runsolver`) that will enforce some limits on both used memory and total CPU time. Solvers will be run inside a sandbox that will prevent unauthorized use of the system (network connections, file creation outside the allowed directory, among others).

Solvers can be run as either 32 bits or 64 bits applications. If you submit an executable, you are required to provide us with an ELF executable (preferably statically linked). Authors submitting solvers in source form will have to specify if it should be compiled in 32 bits or 64 bits mode.

Two executions of a solver with the same parameters and system resources are expected to output the same result in approximately the same time (so that the experiments can be repeated).

## 6.1 Command Line

During the submission process, you will be asked to provide the organizers with a suggested command line that should be used to run your solver. In this command line, you will be asked to use the following placeholders, which will be replaced by the actual information by the evaluation environment.

- `BENCHNAME` will be replaced by the name of the file containing the XCSP3 instance to solve. Obviously, the solver must use this parameter or one of the following variants: `BENCHNAMENOEXT` (name of the file with path but without extension), `BENCHNAMENOPATH` (name of the file without path but with extension), `BENCHNAMENOPATHNOEXT` (name of the file without path nor extension).
- `RANDOMSEED` will be replaced by a random seed which is a number between 0 and 4294967295. This parameter **MUST** be used to initialize the random number generator when the solver uses random numbers. It is recorded by the evaluation environment and will allow to run the program on a given instance under the same conditions if necessary.
- `TIMELIMIT` (or `TIMEOUT`) represents the total CPU time (in seconds) that the solver may use before being killed. May be used to adapt the solver strategy.
- `MEMLIMIT` represents the total amount of memory (in MiB) that the solver may use before being killed. May be used to adapt the solver strategy.
- `NBCORE` will be replaced by the number of processing units that have been allocated to the solver. Note that, depending on the available hardware, a processing unit may be either a processor, a core of a processor or a “logical processor” (in hyper-threading).
- `TMPDIR` is the name of the only directory where the solver is allowed to read/write temporary files
- `DIR` is the name of the directory where the solver files will be stored

Examples of command lines:

```
DIR/mysolver BENCHNAME RANDOMSEED
DIR/mysolver --mem-limit=MEMLIMIT --time-limit=TIMELIMIT --tmpdir=TMPDIR BENCHNAME
java -jar DIR/mysolver.jar -c DIR/mysolver.conf BENCHNAME
```

As an example, these command lines could be expanded by the evaluation environment as:

```
/solver10/mysolver /tmp/zebra.xml 1720968
/solver10/mysolver --mem-limit=900 --time-limit=1200 --tmpdir=/tmp/job12345 /tmp/zebra.xml
java -jar /solver10/mysolver.jar -c /solver10/mysolver.conf /tmp/zebra.xml
```

The command line provided by the submitter is only a suggested command line. Organizers may have to modify this command line (e.g., memory limits of the Java Virtual Machine (JVM) may have to be modified to cope with the actual memory limits).

The solver may also (optionally) use the values of the following environment variables:

- `TIMELIMIT` (or `TIMEOUT`) (the number of seconds it will be allowed to run)
- `MEMLIMIT` (the amount of RAM in MiB available to the solver)

- `TMPDIR` (the absolute pathname of the only directory where the solver is allowed to create temporary files)

After `TIMEOUT` seconds have elapsed, the solver will first receive a `SIGTERM` to give it a chance to output the best solution it found so far (in the case of an optimization problem). One second later, the program will receive a `SIGKILL` signal from the controlling program to terminate the solver.

**The solver cannot write to any file except standard output, standard error and files in the `TMPDIR` directory. A solver is not allowed to open any network connection or launch unexpected external commands. Solvers may use several processes or threads. Children of a solver process are allowed to communicate through any convenient means (Pipes, Unix or Internet sockets, IPC, ...). Any other communication is strictly forbidden. Solvers are not allowed to perform actions that are not directly related to the resolution of the problem.**

## 6.2 Output Format

To communicate their answers, solvers must print messages to the standard output and those messages will be used to check the results. The first two characters of a line allow us to classify it into different categories, which indicate the meaning of the line. With the exception of "o" lines, there is no specific order imposed on the lines output by solvers.

- **status line**

This line starts by the two characters: lower case `s` followed by a space (ASCII code 32). Only one such line is allowed, and it is mandatory. This line gives the answer of the solver. It must be one of the following answers:

- **s UNSUPPORTED**

This line should be printed by the solver when it discovers that the `XCSP3` instance contains a non-supported feature. As an example, a solver that cannot deal with a global constraint should print this line when such a constraint is present.

- **s SATISFIABLE**

This line indicates that the solver has found a solution, and in such a case, a "v" line (see below) is mandatory. For `CSP`, the solver answers `SATISFIABLE` when it has found a solution. For `COP`, the solver answers `SATISFIABLE` when it has found a solution that it couldn't prove to be optimal.

- **s OPTIMUM FOUND**

This line must be printed when the solver has found an optimal solution for a `COP` instance, and in such a case, a "v" line (see below) is mandatory. This answer implies that the solver has proved that no other solution can give a better value of the objective function. This answer must not be used for `CSP` instances.

- **s UNSATISFIABLE**

This line must be output when the solver can prove that the instance has no solution.

- **s UNKNOWN**

This line may be output in any other case, i.e. when the solver is not able to tell anything about the instance.

It is of uttermost importance to respect the exact spelling of these answers. Any mistake in the writing of these lines will cause the answer to be disregarded.

Solvers are not required to provide any specific exit code corresponding to their answer.

If the solver does not output a status line, or if the status line is misspelled, then UNKNOWN will be assumed.

- **values line**

This line starts by the two characters: lower case v followed by a space (ASCII code 32). It is mandatory when the instance is satisfiable. More than one "v " line is allowed but the evaluation environment will act as if their content was merged.

If the solver finds a solution (i.e., if the solver outputs "s SATISFIABLE" or "s OPTIMUM FOUND"), it must provide a solution. For CSP or COP, this solution is an instantiation that satisfies every constraint. For COP, this instantiation must be such that the value of the objective function corresponds to the best one that the solver was able to find.

Solutions must respect the format described in Section 2.11 of [BLAP17]. However, it is important to note that the attributes `type` and `cost` that can be associated with the element `<instantiation>` are **not required** in the context of the competition. These attributes, if present, will simply be ignored.

Importantly, the solution can be output on several successive "v " lines, provided that each "v " line must be terminated by a Line Feed character (the usual Unix line terminator `'\n'`). A "v " line that does not end with that terminator will be ignored because it will be considered that the solver was interrupted before it could print a complete solution.

As an illustration, the following output is valid for the COP instance (Example 4) given in Chapter 1 of [BLAP17]:

```
v <instantiation type="optimum" cost="1700">
v <list> b c </list>
v <values> 2 2 </values>
v </instantiation>
```

and the following output is valid for the CSP instance (Example 25) given in Section 2.11 of [BLAP17]:

```
v <instantiation type="solution">
v <list> x[] </list>
v <values> 1 1 2 * </values>
v </instantiation>
```

As the attributes `type` and `cost` are not required (and simply ignored by our environment), we could have written:

```
v <instantiation>
v <list> b c </list>
v <values> 2 2 </values>
v </instantiation>
```

and

```
v <instantiation>
v <list> x[] </list>
v <values> 1 1 2 * </values>
v </instantiation>
```

- **objective line**

These lines start by the two characters: lower case o followed by a space (ASCII code 32). **These lines are mandatory for incomplete solvers.** As far as complete solvers are concerned, they are not strictly mandatory but solvers are strongly invited to print them. These lines are only relevant for COP instances.

Whenever the solver finds a solution with a better value of the objective function, it is asked to print an "o " line with the current value of the objective function. Therefore, an "o " line must contain the lower case o followed by a space and then by an integer that represents the better value of the objective function. "o " lines should be output as soon as the solver finds a better solution and be ended by a standard Unix end of line character ('\n'). Programmers are advised to flush immediately the output stream.

As an example; let us consider Example 2 in Chapter 1 of [BLAP17]. Let us assume that the solver finds first this solution:

```
<instantiation id='sol1' type='solution' cost='450'>
  <list> b c </list>
  <values> 0 1 </values>
</instantiation>
```

and later:

```
<instantiation id='sol2' type='solution' cost='1700'>
  <list> b c </list>
  <values> 2 2 </values>
</instantiation>
```

which is finally proved to be optimal by the solver. The output by the solver can be (using this time only one "v " line):

```
o 450
o 1700
s OPTIMUM FOUND
v <instantiation> <list> b c </list> <values> 2 2 </values> </instantiation>
```

The evaluation environment will automatically timestamp each of these lines so that it is possible to know when the solver has found a better solution and how good the solution was. The goal is to be able to analyze the way solvers progress toward the best solution. As an illustration, here is a sample of the output of a solver, with each line timestamped (first column, expressed in seconds of wall clock time since the start of the program).

```
0.00    c Time Limit set via TIMEOUT to 1800
0.51    c Initial problem consists of 6774 variables and 100 constraints.
0.55    c preprocess terminated. Elapsed time: 0.45
0.55    c Initial Lower Bound: 0
0.63    o 235947
```

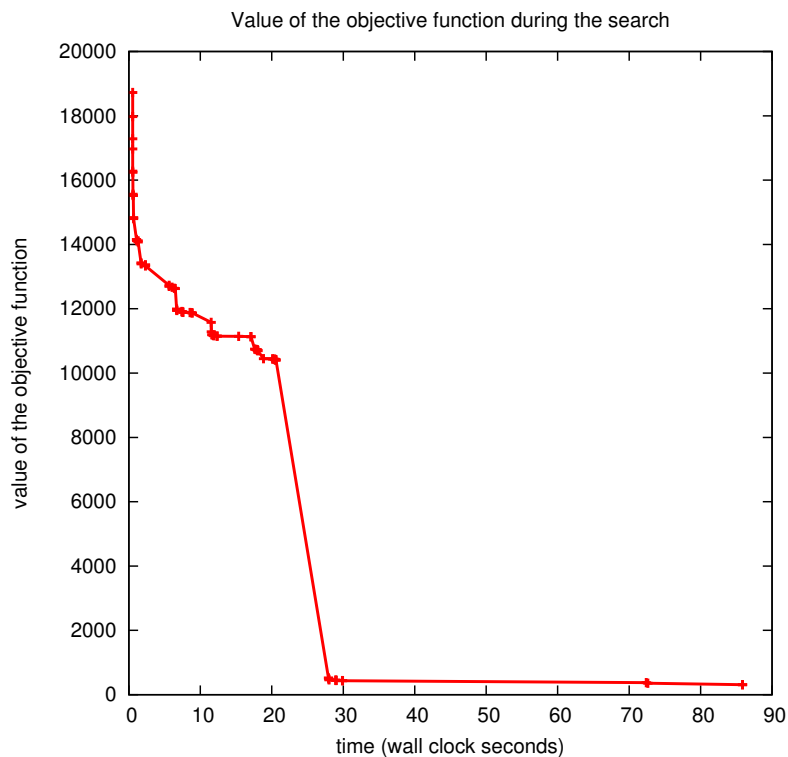


```

0.63      o 226466
0.63      o 217758
0.75      o 186498
1.16      o 178319
2.42      o 168389
3.13      c Restart #1 #Var: 6774 LB: 0 @ 3.03
4.89      c Restart #2 #Var: 6774 LB: 0 @ 4.79
5.73      o 160358
6.44      o 159206
7.52      o 150077
9.09      o 149533
12.14     o 140853
17.74     o 140264
19.61     o 131636
29.81     o 15450
34.00     o 7066
41.66     o 5000
84.01     o 3905
84.01     c NEW SOLUTION FOUND: 3905 @ 83.873
84.61     s OPTIMUM FOUND
84.61     v ... // solution not shown here for space reasons
84.61     c Total time: 84.478 s

```

and here is an example of graph which can be generated from such 'o' lines:



- **diagnostic line**

These lines are optional and start with the two following characters: lower case d followed by a space (ASCII code 32). Then, a keyword followed by a value must be given on this line.

More precisely, a *diagnostic* is a (name,value) pair that gives an information about the work carried out by the solver. As indicated above, each diagnostic is a line of the form

'd NAME value', where NAME is a sequence of letters describing the diagnostic, and value is a sequence of characters defining its value. The following diagnostic is predefined:

**WRONG\_DECISIONS:** The total number of wrong decisions which have been carried out (as defined in [BZF04]).

Contestants wishing to record other diagnostics than the one listed before above should inform the organizers.

- **comments line**

A line which is not one the special lines defined above, or which explicitly starts with the two characters: lower case c followed by a space (ASCII code 32) is a comment line, and is ignored. These lines are thus optional and may appear anywhere in the solver output.

They contain any information that authors want to emphasize, such as #backtracks, #flips,... or internal CPU time. They are recorded by the evaluation environment for later viewing but are otherwise ignored. At most one megabyte of solver output will be recorded. So, if a solver is very verbose, some comments may be lost.

Submitters are advised to avoid printing comment lines which may be useful in an interactive environment but otherwise useless in a batch environment. For example, printing comment lines with the number of constraints read so far only increases the size of the logs with no benefit.

If a solver is really too verbose, the organizers will ask the submitter to remove some comment lines.

## 6.3 Special Considerations for Incomplete Solvers

Complete solvers are solvers which can always decide the satisfiability of a CSP instance and the optimality of a COP instance, provided that enough time and memory are given. Incomplete solvers may loop endlessly in a number of cases; local search algorithms are examples of incomplete solvers. Both kinds of solvers are welcome in this competition. Submitters will have to indicate if their solver is complete or incomplete on the submission form.

### 6.3.1 Complete solvers

There is no special requirement about complete solvers. See the input and output format that all solvers must respect for details.

### 6.3.2 Incomplete solvers

Incomplete solvers are definitely welcome in the competition.

For CSP, an incomplete solver will stop as soon as it finds a solution and will time out if it can't find one. The only difference with a complete solver is that it will time out systematically on unsatisfiable instances.

For COP, an incomplete solver will systematically time out because it will be unable to prove that it has found the optimum solution. Yet, it may have found the optimum value well before the time out. In order to get relevant information in these categories, an incomplete solver must fulfill two requirements:

1. it must intercept the SIGTERM signal sent to the solver on timeout and output either "s UNKNOWN" or "s SATISFIABLE" with the "v " line(s) corresponding to the best solution it has found
2. it MUST output an "o " line whenever it finds a better solution so that, even if the solver always timeouts, the timestamp of the last "o " line indicates when the best solution was found. Keep in mind that it is the evaluation environment which is in charge of timestamping "o " lines.

## 6.4 Special Considerations for Parallel Solvers

The execution environment will bind the solvers to a subset of all available processing units. The environment variable NBCORE will indicate how many processing units have been granted to the solver. The solver will not have access to more processing units than NBCORE. This implies that if the solver uses  $x$  threads or processes (with  $x > \text{NBCORE}$ ),  $x - \text{NBCORE}$  threads or processes will necessarily sleep at one time.

As an example, if the competition is run on hosts with 2 quad-core processors (8 cores in total), several scenarios are possible:

- one single solver is run on the host, it is allowed to use all 8 cores (NBCORE=8).
- two solvers are run simultaneously, each one being assigned to a given processor (which means that a solver is assigned 4 cores, hence NBCORE=4).
- 4 solvers are run simultaneously, each one being assigned to a fixed set of 2 cores (belonging to the same CPU), hence NBCORE=2.
- more generally, a single solver may be assigned any number  $x$  of cores (from 1 to 8 in this example) to simulate the availability of  $x$  processing units.

The solver might use the NBCORE environment variable to adapt itself to the number of available processing units.

A solver must not modify its processor affinity (calls to `sched_setaffinity(2)` or `taskset(1)`) to get access to a processing unit that was not initially allocated to the solver. It may however modify its processor affinity to use a subset of the initially allocated processing units.

## 7 Entering the Competition

Contestants can enter the competition with one or two solvers per track. Contestants are expected to submit their solver(s) and contribute some instances (as many instances as wished). Submitted instances will be made available on the evaluation web site shortly after the actual beginning of the competition. We cannot accept benchmarks which cannot (for various reasons) be publicly available (because anyone must be able to reproduce the experiments of the competition). Each contestant will have the possibility to select 5 instances with the guarantee that they will be used for the competition. In a second stage, they will also have to submit a position paper (at least 2 pages) indicating the main components of the submitted solver(s).

Of course, we expect that contestants propose solvers that recognize XCSP3 (either natively or by embedding a conversion procedure).

The deadline for submitting both benchmarks and solvers is April 29, 2018. Submission of solvers and benchmarks will be possible online in March 2018 at <http://www.cril.fr/XCSP18/>.

## 8 Ranking

Basically, solvers will be ranked on the number of times a solver is able to give the best answer obtained during the competition. Ties will be broken on the cumulated CPU/wall-clock time to give these answers. Other ranking schemes may be introduced to help identify remarkable features.

**Wrong Answers.** Note that a solver is declared to give a wrong answer in the following cases:

- It outputs UNSATISFIABLE for an instance which can be proved to be satisfiable.
- For CSP and COP, it outputs SATISFIABLE or OPTIMUM FOUND, but provides an instantiation that does not satisfy every constraint. The only exception is when the solver outputs an incomplete "v " line (which does not end by '\n') in which case it is assumed that the solver was interrupted before it could output the complete model and the answer will be considered as UNKNOWN.
- It outputs OPTIMUM FOUND but there exists an instantiation with a better value of the objective function/cost than the one corresponding to the printed solution.

**When a solver provides even one single wrong answer in a given track, the solver's results in that track will be excluded from the final evaluation results because they cannot be trusted.** Exceptionally, the organizers may decide to present separately the results of such a solver but only if it obtained particularly good results and if a detailed explanation of the problem as well as a correction is provided by the submitters.

A solver that ends without giving any solution, or just crashes for some reason (internal bugs...), is simply considered as giving an UNKNOWN result.

## 9 Committees

### 9.1 Organization

Olivier Roussel and Christophe Lecoutre from CRIL.

They can be reached at [roussel@cril.fr](mailto:roussel@cril.fr) and [lecoutre@cril.fr](mailto:lecoutre@cril.fr).

### 9.2 Judges

The competition jury is composed of three judges who are in charge of taking decisions when rules are unclear and of validating the results of the competition. The three judges will be announced later.

We recall that anybody is welcome to submit new benchmarks (to the judges), and that each participant has the possibility of selecting exactly 5 instances (while proposing many more instances that may finally be selected or not), and must inform the jury of his/her selection. Submitting instances will be made possible through the competition website.

## References

- [BLAP17] F. Boussemart, C. Lecoutre, G. Audemard, and C. Piette, *XCSP3: An integrated format for benchmarking combinatorial constrained problems*, Tech. Report arXiv:1611.03398, Specifications 3.0.5, CoRR, 2016-2017, Available from <http://www.xcsp.org/format3.pdf>.
- [BZF04] C. Bessiere, B. Zanuttini, and C. Fernandez, *Measuring search trees*, Proceedings of ECAI'04 workshop on Modelling and Solving Problems with Constraints, 2004, pp. 31–40.